

IRQL – Yet Another Language for Querying Semi-Structured Data?

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Abstract

In this paper we describe the basic ideas and concepts behind the Information Retrieval Query Language (IRQL). The principal focus of IRQL development is the integration of concepts of information retrieval and database query languages. Therefore, we will be able to exploit the structure of documents and can additionally use information retrieval techniques regardless of whether the structure is known or not. Our approach develops a query language that is compatible with the recently adopted SQL99 standard and information retrieval clauses (e.g. boolean retrieval). It then integrates features of database query languages such as (1) exploiting the document's structure and (2) restructuring (including linking of multiple documents) and information retrieval techniques such as (I) content-based retrieval, (II) ranking, and (III) relevance feedback. Our data model extends the object-relational model and additionally supports an abstraction of attribute names. We evaluate IRQL queries by mapping them to queries supported by existing systems such as object-relational DBMSs, full-text DBMSs, or conventional search engines, and post processing the results supplied by these systems, if necessary.

IRQL is used as an internal query language for structured and semi-structured data in the GETESS project.

1 Introduction

During the last years the WWW became generally accepted as a medium to publish various kinds of information (documents). In general, this information can be categorized as structured and semi-structured/unstructured. Although storing and querying of structured data (e.g. using relational DBMSs) are well understood, there is still no agreement in managing semi-structured data (e.g. data kept in files; possibly using XML). Keeping this potential heterogeneity in mind, it is quite difficult to search for particular information. On the one hand, there are many search engines (e.g. Altavista or Infoseek) that permit the search for particular documents as it relates to their content, but these search engines are often not capable of exploiting

the structure of documents in order to support advanced queries. Additionally, often data stored in DBMSs are not taken into account, although these search engines could benefit from the features of database query languages. On the other hand, pure database query languages are also inappropriate for querying heterogeneous semi-structured data as it relates to documents. These query languages certainly support operations on structured parts of documents, but the ability to query semi-structured data is rather limited and often realized by vendor-specific extensions to the DBMS.

In this paper we describe the basic concepts of an Information Retrieval Query Language (IRQL). The principal focus of IRQL development is to unite the features of database query languages such as (a) access to the data’s structure, (b) use of type specific information, (c) restructuring, and (d) linking of data and information retrieval techniques such as (i) content-based retrieval, (ii) vague queries, (iii) ranking, and (iv) relevance feedback into a single query language. Thus, IRQL allows us to query both structured and heterogeneous semi-structured data related to documents.

Our approach is to store our data in existing systems such as object-relational DBMSs, relational DBMSs, or full-text DBMSs. Therefore, we implement IRQL on top of these systems. In principle, we evaluate IRQL queries by mapping them to the query languages supported by the corresponding platform as illustrated in Figure 1. Obviously, we have to post process the results delivered by these platforms as none of the systems support all of the IRQL features. This post-processing is either done by wrappers or compensators. Essentially, the difference between wrappers and compensators is that compensators are “big” wrappers, i.e. compensators encapsulate systems that support only a very limited set of the features of IRQL. To be more concrete, assume that we want to evaluate an IRQL query on top of a relational database system and the query includes selection, projection, proximity search, and ranking functions. First, we map the query’s relational parts (e.g. selection and projection) of this query to the SQL supported by the RDBMS and evaluate the query. As a result, we obtain a superset of the “real” result. Now we ask the compensator to compute the answer to the IRQL query using the result from the RDBMS and the parts of the IRQL query that could not be mapped to the RDBMS query language (e.g. proximity search and ranking).

IRQL is used in the GETESS project [SBB⁺99a, SBB⁺99b] to query the summaries of linguistically analyzed web documents.

The remainder of this paper is organized as follows: In Section 2, we discuss some related work and compare other approaches with IRQL. We describe our data model in Section 3. Section 4 presents the syntax of IRQL using some examples. In Section 5, we present two applications of IRQL. We conclude with a summary in Section 6 and mention some future works.

2 Related Work

IRQL unites concepts from database query languages, query languages for semi-structured data, and information retrieval. In this section, we discuss some query language proposals that are related to these areas (i.e. we focus primarily on query

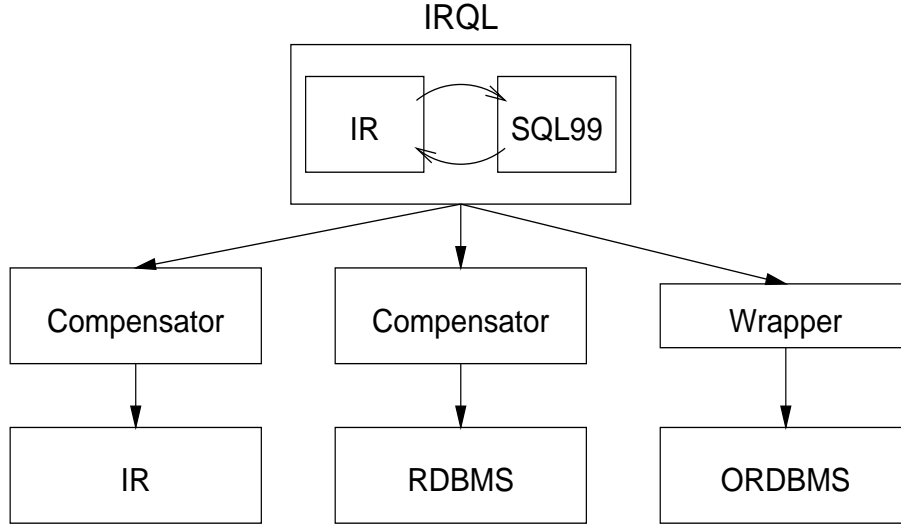


Figure 1: implementation of IRQL

languages for semi-structured and web data) and compare them with our approach. We do not discuss XML query languages here because these query languages don't consider the integration of information retrieval techniques and, in principle, XML data can also be queried using some of the following query languages.

Most of the existing search engines (e.g. Altavista or Infoseek) use information retrieval techniques to search for particular (web) documents. Users describe their search criteria by entering keywords, phrases, or combinations using boolean operators. However, these search engines don't normally take the document's structure into account, and only selections are primarily supported. Features that are typical for DBMSs (and also for IRQL) like restructuring (e.g. projection) or joins are still missing.

Access to the document structure is supported by freeWAIS-sf [PFH95, Pfe95]. Documents can be partitioned into a set of attribute-value pairs. The set of possible attributes is defined by the document type. One of the pre-defined types is HTML and further types can be defined by the user. Queries are also limited to selections, but parts of the document can be queried via attribute names. Users can describe their search criteria using free text, phrases, wildcards, soundex and proximity expressions, as well as combinations of these using boolean operators. Comparisons of numeric values are supported, too. Besides the access to the document's structure in IRQL, we also allow for the restructuring of data.

In the past, the disadvantages of exclusively using information retrieval techniques to query semi-structured data [Abi97, Clu97] has been pointed out by several authors. As a result, there are numerous proposals and implementations that also integrate concepts of database query languages. A survey can be found in [FLM98].

Lorel [AQM⁺97] is the query language of the Lore system [MAG⁺97]. Syntactically, Lorel is based on OQL. Semi-structured data are supported by using OEM

graphs as a data model and by extending OQL with appropriate features. These extensions include (a) implicit type casts (type coercion) and (b) regular path expressions. Path expressions and path variables support queries on unknown or partially known schemas and on the schema itself. Implicit type casts, called type coercion, address the heterogeneity of semi-structured data. Some variants of content-based retrieval (e.g. soundex search) are provided by corresponding predicates. The drawback of this data model is the missing support of ordered collections¹. As a consequence, for example, no ranking criteria can be specified at the language level.

WebSQL [AMM97] is based on the relational model and supports an SQL-like query language. Additional features of WebSQL include dynamic creation of extensions based on content and link structure of web documents, and path expressions. HTML tags are treated as attribute names in order to access the structure of web data. The restructuring of HTML pages is not supported.

Compared with OEM graphs, WebOQL [AM98] uses an improved data model. Hypertrees facilitate the modelling of nested structures and further support ordered collections. Using the “web” as a data type is the key to providing a number of operations for restructuring data. The query language is based on OQL and provides some further possibilities, e.g. the creation of query results. Content-based retrieval is supported by a grep operator. In our approach, we support further means of information retrieval, such as term weighting and ranking.

UnQL [BDHS96] uses a graph-based data model. The query language supports operations such as selection, projection, join, and grouping, as well as path expressions. Both ordered collections and content-based retrieval are not supported.

W3QL [KS95] is the SQL-like query language of W3QS. The focus of this query language’s development is the reuse of available tools. For example, predicates that realize content-based retrieval are implemented using external tools. Both nesting of queries and restructuring of data are not supported.

The aim of the development of the Strudel query language StruQL [FFLS97] is to provide means for restructuring existing data. In StruQL, semi-structured data are modelled as an OEM graph. The supported query operations include navigation using path expressions, projection, and selection as well as operations for restructuring of existing graphs and for creation of new graphs. In principle, user-defined predicates could be used to implement content-based retrieval.

WebLog [LSS96] is based on SchemaLog and supports access to the structure of documents and content-based retrieval by using built-in or user-defined predicates. The restructuring of data is supported, too. As in IRQL, it is possible to express recursive queries.

In WQL [LSCH98], both the web and the structure of the individual documents are modelled. The query language implements projection, selection, sorting, and grouping. Content-based retrieval and querying the structure of web documents are supported. From our point of view, missing features are dynamically creating extensions, nesting, and restructuring.

¹Recently, Lorel’s data model has been extended to support XML data [GMW99]. Therefore, ordered subelements can now be modelled. To the best of our knowledge, it is not possible to express user-defined rankings in Lorel like it is in IRQL.

Apart from the query languages for semi-structured data mentioned in this section, there are also proposals to extend DBMSs. For example, the SQL/MM proposal [SQL95] defines a data type *full text* whose operations support content-based retrieval for those data that are stored using this data type. Different implementations are made available by commercial companies in the form of text extenders, data blades, and so on.

Figure 2 summarizes the features of the query languages mentioned in this section. We use the following notation to assess the features: + supported, \pm supported with limitations, \mp supported to some extent, and – not supported.

3 Data Model

The principal focus of IRQL development is to implement a query language that allows to query both structured data (this implies that IRQL adopts features from database query languages) and semi-structured data. Therefore, IRQL also includes information retrieval techniques. Apart from the separate use of these two query types, IRQL integrates both possibilities.

Our data model extends the object-relational data model. There are atomic types such as² `integer`, `float`, `boolean`, and `string`. Composite types include the collections `set`, `bag`, `list`, and `array` as well as the `struct` constructor. In order to model heterogeneous semi-structured data, we introduce a composite type `doc` similar to the `struct` constructor, but querying data of this type does not produce any type checking errors. Within a `doc` type, we also allow for referencing non-existent labels.

We illustrate the `doc` type by an example. Figure 3 shows a structured instance of type

```
set<struct(name:string, place:string, rooms:integer)>.
```

Figure 4 shows a semi-structured instance containing two tuples. The first tuple type could be

```
struct(name: string, equipment: set<string>,
       drinks: set<string>,
       price: struct(single: integer, double: integer))
```

while the second tuple could be of type

```
struct(name: string, equipment: set<string>,
       price: struct(single: integer, double: integer,
                     twin: integer, app: integer)
       cards: set<string>).
```

In our approach, heterogeneous data is modelled as `doc` type. Therefore, the tuples from Figure 4 are typed as

```
doc(name: string, equipment: set<string>,
    drinks: set<string>,
    price: doc(single: integer, double: integer))
```

²We plan to support a large subset of the atomic types mentioned in [SQL99a].

	Lorel		WebOQL		WebSQL		UnQL		W3QL	
	OEM graph	hyper tree	relational	labeled graph	labeled graph	SQL	structural recursion	SQL	labeled graph	SQL
data model	OQL	OQL	SQL	SQL	SQL	SQL	SQL	SQL	SQL	SQL
language style	+	+	+	+	+	+	+	+	+	+
path expressions	+	-	-	-	-	-	-	-	-	-
type coercion	+	-	-	-	-	-	-	-	-	-
updates	-	-	-	-	-	-	-	-	-	-
strong typing	-	-	-	-	-	-	-	-	-	-
recursion	-	-	-	-	-	-	-	-	-	-
content bases retr.	±	±	±	±	±	±	±	±	±	±
ranking	-	-	-	-	-	-	-	-	-	-
term weighting	-	-	-	-	-	-	-	-	-	-
ordered collections	-	+	-	-	-	-	-	-	-	-
linking different docs.	+	+	+	+	+	+	+	+	+	+
querying the schema	+	-	-	-	-	-	+	-	-	-

	StruQL		WebDB		WebLog		freeWAIS-sf		IRQL	
	labeled graph	object-relational	relational	relational	relational	relational	relational	relational	relational	relational
	Datalog	SQL	Datalog	SQL	Datalog	SQL	Datalog	SQL	Datalog	SQL
data model	+	-	-	-	-	-	-	-	-	-
language style	-	-	-	-	-	-	-	-	-	-
path expressions	-	-	-	-	-	-	-	-	-	-
type coercion	-	-	-	-	-	-	-	-	-	-
updates	-	-	-	-	-	-	-	-	-	-
strong typing	+	-	-	-	-	-	-	-	-	-
rekursion	-	±	±	±	±	±	±	±	±	±
content based retr.	-	-	-	-	-	-	-	-	-	-
ranking	-	-	-	-	-	-	-	-	-	-
term weighting	-	-	-	-	-	-	-	-	-	-
ordered collections	-	+	-	+	-	+	-	+	-	+
linking different docs.	+	+	+	+	+	+	+	+	+	+
querying the schema	-	-	-	-	-	-	-	-	-	-

Figure 2: features of some query languages

<i>name</i>	<i>place</i>	<i>rooms</i>
Neptun	Warnemünde	⊥
Hübner	Warnemünde	95
Mecklenburger Hof	Rostock	21
Atrium Hotel Krüger	Sievershagen	59

Figure 3: structured instance (R1)

<i>name</i>	<i>equipment</i>	<i>drinks</i>	<i>price</i>	
			<i>single</i>	<i>double</i>
Hübner	bathroom shower WC ...	beer wine	195	235

<i>name</i>	<i>equipment</i>	<i>price</i>				<i>cards</i>
		<i>single</i>	<i>double</i>	<i>twin</i>	<i>app.</i>	
Krüger	bathroom shower WC phone ...	95	150	185	135	American Express Visa Diners Club

Figure 4: heterogeneous instance (R2)

and

```
doc(name: string, equipment: set<string>,
    price: doc(single: integer, double: integer,
               twin: integer, app: integer)
    cards: set<string>).
```

As we cannot unify two semi-structured types in general, but need some kind of notation for semi-structured types, we omit the attribute-value pairs from the `doc` type and instead use `doc` for both types. Therefore, the instance from Figure 4 could be typed as `set<doc>`. Instances of the `doc` type are only subject to limited type checks so that, for example, $\pi_{cards}(R2)$ only delivers the credit cards accepted by the Krüger hotel.

While the modelling of semi-structured heterogeneous data as previously discussed is not new, our main contribution concerning the data model is to allow a set of attributes to be abstractly referenced by single attribute names as illustrated in Figure 5.

For example, we introduce two default attributes if the corresponding data originated in web documents: *source* indicates the document’s URL and *complete_content*

document	source	complete_content						
		metadata			text			
	article	authors	title	year	abstract	sections	appendix	references
	http://e-lib.informatik.uni-rostock.de/2000/DBIS/IRQL/bncod.ps	Andreas Heuer Denny Priebe	IRQL-Yet another language for querying semistructured data?	2000	In this paper we describe the basic ideas and concepts behind the ...	During the last years the WWW became generally accepted as a medium to publish various kinds of information (documents). ...		[ABI97] Serge Abiteboul. Querying Semi-Structured Data. In Foto N. Afrati and Phokion Kolatis, editors, Database Theory - ICDT'97 ...

Figure 5: abstract attributes

indicates the full text of the original page. As Figure 5 shows, *complete_content* is an abstraction of a set of different attributes, e.g. *metadata* such as *authors* and *text*. *Text*, in turn, is another abstraction of further attributes such as the abstract or the references of the modelled article.

In contrast to object-oriented or object-relational database models, the attributes *complete_content* and *text* are no tuple-valued attributes. For example, *complete_content* would consist of two different components *metadata* and *text* in the object models. Here, *complete_content* is considered as one text value again. The advantage of this kind of abstraction operator is the usability for information retrieval operations. If useful, the *complete_content* value can be seen as one atomic value.

The problem that object-oriented and object-relational models (that are used as implementation models) do not support this kind of abstraction is hidden from the user: Our abstraction operator is implemented on top of existing object-oriented and object-relational concepts.

4 Language

The aim of IRQL development is to integrate database query languages and information retrieval techniques. Similiar to Lorel, our approach is to realize a query language in the style of SQL, but we additionally support information retrieval techniques by adding new clauses. Like some of the query languages mentioned in Section 2, we also change the type checking rules of SQL to also support querying semi-structured heterogeneous data.

Because of these demands, we take the recently adopted SQL99 standard [SQL99a, SQL99b] as a starting point. We plan to implement a large subset of the proposed syntax in order to be able to answer queries conforming to this standard. Using examples, we subsequently show possibilities for querying semi-structured data and integrating information retrieval techniques into IRQL.

4.1 Structured and Semi-structured Data

The data model described in Section 3 supports querying structured and semi-structured data. Structured composite data are modelled as elements of the `struct` data type and are therefore subject to the strong type checking as found in e.g. SQL99. Semi-structured data are modelled as elements of a special data type (`doc`). We modify the type checking rules for instances of this data type³ so that meaningful queries are possible, even if the schema is not known or only partially known. These modifications include: (1) incompatible data types are casted to compatible types, if necessary and (2) in heterogeneous data, non-existent attributes may be referenced. The concrete semantics are dependent on the type operation used. For example, a non-existent attribute referenced in a projection is ignored for each tuple it does not appear in. Therefore, the operation’s result is again heterogeneous. Selection predicates referencing such attributes are evaluated to false. (3) Partially known schemata can be queried using path expressions and path variables. For example, the query

```
select r.name, ##z, r.cards
from   R2 r, r.price.{.*}z
where  ##z < 200
```

		<i>name</i>	<i>single</i>	<i>double</i>	
		Hübner	195	235	
<i>name</i>	<i>single</i>	<i>double</i>	<i>twin</i>	<i>app.</i>	<i>cards</i>
Krüger	95	150	185	135	American Express Visa Diners Club

Figure 6: heterogeneous result

results in the heterogeneous instance shown in Figure 6. `R2` denotes the instance from Figure 4. The elements of this set are typed as `doc` and are therefore subject to the type checking rules mentioned earlier. In the `from` clause, regular path expressions are used and the path variable `z` is declared. The regular path expression `{.*}` expands to all existing attributes below `price`. The appended (optional) label `z` denotes the corresponding variable name. The expression `##z` dereferences the path variable `z` and is substituted by the complete paths (in this example `R2.price.single`, `R2.price.double` for the first tuple, as well as `R2.price.single`, `R2.price.double`, `R2.price.twin` and `R2.price.app` for the second tuple). The attribute `cards` referenced in the `select` clause is ignored while processing the first tuple as there is no such attribute there. If one or more prices are string types, these prices would have to be converted to numeric values to evaluate the predicate `##z < 200`. If such a conversion is not possible, the predicate is evaluated to false.

³Essentially, we adopt the techniques (primarily Lorel’s) used in existing query languages for querying semi-structured data.

<i>name</i>	<i>price</i>	
	<i>single</i>	<i>double</i>
Hübner	195	235
Krüger	95	150

Figure 7: instance (R3)

The integration of structured and semi-structured data is realized by modelling these data as instances of different data types. For example, let the instance shown in Figure 7 be of type

```
set⟨doc(name: string,
        price: struct(single: integer, double: integer))⟩.
```

The query

```
select r.price.single, r.address
from   R3 r
```

delivers the prices of all single rooms because the non-existent attribute address is directly ignored within a doc type. But the query

```
select r.price.app, r.address
from   R3 r
```

leads to a runtime error because the price is modelled as a struct type where non-existent attributes may not be referenced.

As illustrated previously, the difference between operations on structured and semi-structured data can be modified in order to be valid for other operations, too.

4.2 Extensions

In IRQL, there is one basic extension describing all known documents. We call this extension *d_world*. In order to avoid considering all these documents in every query, we introduce some possibilities to create further (e.g. smaller) extensions. Useful criteria include (1) information about how or whether a document can be reached via a particular path, (2) the language of documents, (3) the possibly named document types, and (4) the document's domain. We express each of these as an extension to the **from** clause. The extension (*coll*) of documents that are reachable starting from a given URL is determined by

```
⟨coll⟩ REACHABLE FROM ⟨URL⟩ [DEPTH ⟨value⟩] [LOCAL].
```

As options, a maximum path length (DEPTH parameter) can be specified or only local documents can be chosen. All documents in a given language can be determined by

```
⟨coll⟩ IN LANGUAGE ⟨lang⟩.
```

The clause

`<coll> OF [NAMED] TYPE <tycon>`

creates a document extension for other possibly named type. If the keyword **NAMED** is omitted, *tycon* stands for a type constructor, otherwise it is a label like postscript (PS). Finally,

`<coll> OF DOMAIN <domain>`

determines the extension of all documents of a given domain (e.g. tourism).

4.3 Information Retrieval

In the following, we describe a further extension to SQL99; namely predicates that implement information retrieval techniques (e.g. content-based retrieval, soundex and proximity search, term weighting and ranking of query results). These possibilities are particularly suited for, but not limited to, semi-structured data.

4.3.1 Content-based Retrieval

We denote content-based retrieval by the clause

```
<attribute> CONTAINS <text>
[ATLEAST <value>] [ATMOST <value>]
[WITH WEIGHT <value>]
[CASE SENSITIVE] [SUBSTRING]
[( <value> | NO) ERRORS].
```

The following optional parameters exist: (1) **ATLEAST**, **ATMOST** specifies how often *text* must occur in *attribute*. If the number of occurrences of *text* is not within the specified bounds, the predicate is evaluated to false. If one or both of the parameters are omitted, no limit is assumed. (2) **WITH WEIGHT** specifies the weight of the query term. The default value is a weight of one. (3) By default, the search is case insensitive. This can be changed by specifying the **CASE SENSITIVE** parameter. (4) **SUBSTRING** specifies not only matching word bounds (e.g. spaces) but also searching for any occurrence of the given substring. (5) There is also a possibility for considering typing errors (see glimpse [Har] how this can be realized) by specifying a value for the **ERRORS** parameter. By default, no typing errors are considered. The values of *text* can be either keywords or phrases. Furthermore, we support regular expressions (wildcards) here.

4.3.2 Soundex

The soundex algorithm allows the search for phonetically similar keywords or phrases. We denote the soundex search by

```
<attribute> SOUNDEX <text>
[ATLEAST <value>] [ATMOST <value>].
```

The meaning of **ATLEAST** and **ATMOST** can be taken from Section 4.3.1. Further parameters mentioned there are not meaningful within the context of a soundex search.

4.3.3 Proximity

The next supported concept of content-based retrieval is the proximity search. Using a proximity search, it is possible to specify the distance between two keywords or phrases. The denotation is as follows:

```
<attribute> CONTAINS  
<text> [WITH WEIGHT <value>]  
[<value> <unit>] BEFORE | AFTER <text> [WITH WEIGHT <value>]  
[ATLEAST <value>] [ATMOST <value>]  
[CASE SENSITIVE] [SUBSTRING]  
[[<value> | NO] ERRORS].
```

Here, we only describe the new parameters. The others can be found in Section 4.3.1. The new parameter *unit* can be substituted by a type-dependent unit. For example, if *d* is a L^AT_EX document and a method exists to split this document into sections, then

```
d CONTAINS ‘‘related work’’ 2 SECTIONS BEFORE  
‘‘conclusion’’
```

is a valid predicate that checks whether *d* contains the phrase “related word” not more than 2 sections before the keyword “conclusion”.

4.3.4 Ranking

We support ranking results by user-defined criteria. Syntactically, this is denoted by

```
RANK BY  $f_0, \dots, f_n$   
[LIMIT TO <value>]
```

The f_i denote user-defined functions that define the calculation of the retrieval status value (RSV). The RSV is an attribute that is appended by the **rank by** clause and, after calculating this value, the result is sorted by RSV. Although we next plan to support the vector space model, we don’t need to change our syntax if we implement a probabilistic model, as the following example demonstrates:

```
SELECT RSV, name  
FROM hotels d  
RANK BY d.stars=5, d.beachdist=0
```

In this query we define a ranking using boolean predicates. These predicates are not evaluated to true or false, but define the “best” hotel. Thus, the retrieval status value of one is assigned to a five-stars-hotel directly situated at the beach. Using probabilistic methods, the other hotels are ranked accordingly.

The optional part of the **rank by** clause is used to limit the number of returned elements to *value*. By default, the number of elements is unlimited.

4.3.5 Compatibility with DBQLs and Information Retrieval

On the one hand, compatibility with SQL is achieved if there are no semi-structured data, and therefore, no `doc` type data in any of the extensions queried. In this case, any query that is a valid query within the supported subset of SQL99 is also a valid IRQL query and delivers the same result. On the other hand, compatibility with information retrieval expressions is achieved by transparently mapping these expressions to IRQL queries, as the following example demonstrates: Assume we are interested in a hotel near the beach. Using one of the search engines, we would probably enter

hotel and beach.

This expression is also accepted by IRQL and transparently mapped to⁴

```
select <default_projection>
from   <default_extension>
where  <default_attribute> CONTAINS "hotel" AND
      <default_attribute> CONTAINS "beach".
```

The default values can be adjusted within the IRQL shell. A good choice would be to use `source,title` as `default_projection`, `d_world` as `default_extension`, and `complete_content` as `default_attribute`. In this simple example the resulting query

```
select source,title
from   d_world
where  complete_content CONTAINS "hotel" AND
      complete_content CONTAINS "beach"
```

delivers the expected information.

5 Applications

The IRQL language is used in two different projects:

In the BlueView project⁵, digital library services are developed and partially implemented based on the architecture of virtual document servers. Using standard tools like full-text database or information retrieval systems, object-relational database management systems, and replication and caching services, different heterogeneous local document servers have been integrated into one local server. IRQL is the query language for this integrated local document server because it can be implemented on top of these different platforms.

In the GETESS project⁶, we analyse web documents using linguistic and domain-specific knowledge [SBB⁺99a, SBB⁺99b] and use the data gathered in this way to answer user queries. Here IRQL serves as an internal query language. Within this context, we use some default attributes if the corresponding data originated in web

⁴For simplicity, we ignore the ranking.

⁵<http://wwwdb.informatik.uni-rostock.de/blueview>

⁶<http://www.getess.de>

documents: e.g. the attribute “source” points to the URL where the data are gathered from; the attribute “complete_content” contains the full text of the original page.

6 Conclusion and Future Work

In this paper we present our approach for developing an Information Retrieval Query Language (IRQL). The used data model distinguishes structured and semi-structured heterogeneous data based on type information and supports an abstraction of attribute names. IRQL integrates concepts of database query languages, query languages for semi-structured data, and information retrieval techniques. The starting point of IRQL development is SQL99, which we extend with new clauses to integrate information retrieval techniques. Furthermore, we modify the type system to support semi-structured heterogeneous data. IRQL⁷ is built on top of existing systems such as object-relational DBMSs, relational DBMSs, or full-text DBMSs. The current prototype implementation has been built on top of DB2 and its text extender.

To the best of our knowledge, there is no similar proposal that attempts to integrate features of these three areas.

Future works include the complete formalization of the query language and the development of an algebra.

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